

Diskrete Mathematik HS2025 — Prof. Dennis HOFHEINZ

Marian DIETZ — Milan GONZALEZ-THAUVIN — Zoé REINKE

Exercise sheet 13

This is the exercise sheet number 13, last one of the semester. The difficulty of the questions and exercises are rated from very easy (\star) to hard ($\star\star\star\star$). There is no bonus exercise this week. Because this is the last sheet, some exercises need content that will be covered next Monday and Wednesday.

Exercise 13.1 Warm-Up (\star)

1. Are the formulas $F = P(x) \wedge \neg Q(y)$ and $G = \neg Q(x) \wedge P(y)$ equivalent?
2. Find a formula H which has both x and y as free variables, such that $\forall x H \equiv \forall y H$.
3. What does it mean for a (propositional) calculus to be **complete**?

Exercise 13.2 Tautologies ($\star\star$)

Let F be a formula with free variables x_1, \dots, x_n (and no others). Show that F is valid if and only if $\forall x_1 \dots \forall x_n F$ is valid.

Exercise 13.3 Prenex Normal Form (\star)

For each of the following formulas, find an equivalent formula in the prenex normal form.

1. $(\forall x P(x)) \rightarrow Q(x)$
2. $\forall z \exists y (P(x, g(y), z) \vee \neg \forall x Q(x)) \wedge \neg \forall z \exists x \neg R(f(x, z), z)$

Exercise 13.4 Formulas and Statements ($\star\star$)

For each of the following expressions, determine whether it is syntactically correct, and, if so, whether it is a formula or a statement about formulas¹. If an expression is a statement, decide whether it is true or false (each time justify your answer).

1. $\forall x \exists y (P(z) \leftrightarrow Q(f(f(x, z), y)))$
2. $(\forall x P(x)) \models P(x)$
3. $(P(x) \models P(x)) \equiv Q(x)$
4. $\{P(x), P(f(a))\} \models P(a)$

¹Whenever parentheses are not necessary, they can be omitted. Parentheses do not influence correctness.

Exercise 13.5 The Barber of Zurich (★)

Use Corollary 6.13 to show that there does not exist a barber in Zurich who shaves all those and exactly those who do not shave themselves.

Exercise 13.6 Calculi

1. (★) Decide which of the following rules are correct (= sound) and justify your answers:

$$\begin{array}{lll} \{F\} \vdash_{R_1} F \vee G & \{F \wedge G\} \vdash_{R_2} F & \{\neg(F \wedge G)\} \vdash_{R_3} \neg F \wedge \neg G \\ \{F, F \rightarrow G\} \vdash_{R_4} G & \{F \rightarrow G\} \vdash_{R_5} \neg F \rightarrow \neg G & \{F, G\} \vdash_{R_6} F \wedge G \end{array}$$

2. (★ ★) Let K be the calculus consisting of the **correct** rules of question 1 only. Using K , derive formally the formula $((A \wedge B) \wedge C) \wedge D$ from the following set of formulas:

$$\{(D \wedge A) \rightarrow C, A \wedge B, B \wedge A, (B \vee C) \rightarrow D\}$$

3. (★ ★) Is $K' = \{R_2, R_4\}$ complete? Justify your answer.
4. (★ ★) Give an example of a calculus, which is complete but not sound.

Exercise 13.7 Resolution

1. (★) Prove the following statements using the resolution calculus.
- i) $F = (A \vee B) \wedge (\neg E) \wedge (\neg B \vee D) \wedge (\neg D \vee E) \wedge (\neg A \vee B)$ is not satisfiable.
 - ii) $G = (\neg B \wedge \neg C \wedge D) \vee (\neg B \wedge \neg D) \vee (C \wedge D) \vee B$ is a tautology.
 - iii) $H = A \wedge C$ is a logical consequence of $M = \{A \rightarrow C, B \rightarrow A, A \vee B\}$.
2. (★ ★ ★) Let \mathcal{K} be a finite set of finite clauses and let $\mathcal{K}_0, \mathcal{K}_1, \dots$ be a sequence of applications of derivation rules, such that $\mathcal{K}_0 = \mathcal{K}$ and $\mathcal{K}_i = \mathcal{K}_{i-1} \cup \{K\}$ for all $i > 0$, where $\{K', K''\} \vdash_{\text{res}} K$ for some $K', K'' \in \mathcal{K}_{i-1}$. Show that there exists an n such that $\mathcal{K}_m = \mathcal{K}_n$ for all $m > n$.
Hint: Intuitively, the goal of the question is to show that from a finite set of finite clauses, after a finite number of applications of derivation rules, no new clauses can be derived.
3. (★ ★ ★) Show that the statement from question 2 is no longer true for an infinite set \mathcal{K} of finite clauses. More precisely, let $\mathcal{K} = \{\{A_j, \neg A_{j+1}\} \mid j \in \mathbb{N}\}$. Show that there exists an infinite sequence $\mathcal{K}_0, \mathcal{K}_1, \dots$, such that $\mathcal{K}_0 = \mathcal{K}$ and $\mathcal{K}_i = \mathcal{K}_{i-1} \cup \{K\}$ for all $i > 0$, where $\{K', K''\} \vdash_{\text{res}} K$ for some $K', K'' \in \mathcal{K}_{i-1}$, and for all $i > 0$, $\mathcal{K}_i \neq \mathcal{K}_{i-1}$.

Exercise 13.8 A special calculus (exam FS 2025) (★ ★)

This exercise is taken from the spring exam of 2025.

Consider the calculus consisting of the following derivation rules:

$$\begin{aligned} \emptyset &\vdash_{R_1} F \rightarrow (G \rightarrow F), \\ \emptyset &\vdash_{R_2} (F \rightarrow (G \rightarrow H)) \rightarrow ((F \rightarrow G) \rightarrow (F \rightarrow H)), \\ \{F, F \rightarrow G\} &\vdash_{R_3} G. \end{aligned}$$

Formally derive $A \rightarrow C$ from $\{A \rightarrow B, B \rightarrow C\}$ in the calculus.

No exercise will be graded.